# Solutions Exam Imperative Programming

Tuesday 10 November 2015, 18:30-21:30h

## **Problem 1: Assignments** (20 points)

For each of the following annotations determine which choice fits on the empty line (....). The variables x, y and z are of type int. Note that X, Y and Z (uppercase!) are specification-constants (so not program variables).

```
1.1 /* 2*y + x > 12 */
                                       1.4 / * z + y > 4 */
                                           /* z + 1 + y > 5 */
     /* 2*y + x - 1 > 11 */
                                           x = z + 1;
    x = 2 * y + x - 1;
     /* x > 11 */
                                           /* x + y > 5 */
                                           y = x + y;
     /* x > 11 */ So, answer C.
                                            /* y > 5 */ So, answer C.
1.2 /* 6*x + 3*y + z == 3*X - 2*z */ 1.5 /* x == X, y == Y, z == Z */
    /* 3*(2*x + y + z) == 3*x */
                                           /* x == X, y == Y, x + y + z == X + Y + Z */
    /* 2*x + y + z == x */
                                            z = x + y + z;
                                            /\star x == X, y == Y, z == X + Y + Z \star/
    x = 2 * x + y;
    /* x + z == x */
                                            /* z - x - y == Z, y == Y, z == X + Y + Z */
    z = z + x;
                                           x = z - x - y;
                                           /\star x == Z, y == Y, z == X + Y + Z \star/
    /* z == X */
                  So, answer A.
                                            /* x == Z, z - x == X + Y, z == X + Y + Z */
                                           y = z - x;
                                            /\star x == Z, y == X + Y, z == X + Y + Z \star/ So, answer B.
1.3 / * 7 * X - 3 \le x < 7 * X + 4 * /
                                      1.6 /* x == 3*X - 2*Y + 1, y == X */
                                           /* x == 3*X - 2*Y + 1, 3*y - x + 1 == 2*Y */
    /* 7*X \le x + 3 < 7*(X+1) */
    /* X \le (x + 3)/7 < X + 1 */
                                           /* x == 3*X - 2*Y + 1, (3*Y - X + 1)/2 == Y */
   x = (x + 3)/7;
                                           y = (3*y - x + 1)/2;
   /* X \le x \le X + 1 */
                                           /* x == 3*X - 2*Y + 1, y == Y */
                                           /* (x + 2*y -1)/3 == X, y == Y */
    /* x == X */ So, answer B.
                                           x = (x + 2 * y - 1)/3;
                                            /* x == X, y == Y */ So, answer A.
```

### **Problem 2: Time complexity** (20 points)

In this problem the specification constant N is a non-zero natural number (i.e. N>0). Determine for each of the following program fragments the sharpest upper limit for the number of calculation steps that the fragment performs in terms of N.

```
1. int i = 0, s = 0; 3. int i = 0, j = N, s = 0; 5. int i = 0, s = 1;
  while (s < N) {
                                while (j-i > 1) {
                                                             for (i=0; i < N; i+=2) {
    s += i;
                                  int k = (i+j)/2;
                                                                s = 2*s;
    i++;
                                  if (i%2 == 0) {
                                    i = k;
                                                              O(N), answer C.
  O(\sqrt{N}), answer B.
                                  } else {
                                    j = k;
                                                            6. int i, j, s = 0;
                                  }
2. int i, j, s = 0;
                                                              for (i=0; i < N; i++) {
                                  s++;
  for (i=0; i < N; i++) {
                                                                for (j = i*i; j > 0; j = j/2) {
    for (j=N; j > i; j--) {
                                O(\log N), answer A.
                              4. int i, j=0, s=0;
                                                              O(N \log N), answer D.
  O(N^2), answer E.
                                for (i=0; i < N; i++) {
                                  j += i;
                                while (j > 0) {
                                  s += j;
                                  j--;
                                O(N^2), answer E.
```

#### **Problem 3: bit reversal of integers** (10 points)

The binary representation of the number 23 is 10111, since it is equal to  $1 \times 2^4 + 0 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 = 16 + 4 + 2 + 1$ . If we reverse the bit string 10111 we get 11101 which represents the number  $29 = 1 \times 2^4 + 1 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 = 16 + 8 + 4 + 1$ . Therefore, we call 29 the *bit reversal* of 23. Note that integers on a PC are represented by 32 bits, so the binary representation of 23 on a PC consists actually of 27 leading zeros followed by 10111. In this problem we omit all leading zeros. Therefore, the bit reversal of the number 1 is simply 1 itself.

```
#include <stdio.h>
#include <stdlib.h>

int bitrev(int n) {
    int m = 0;
    while (n > 0) {
        m = 2*m + (n%2);
        n /= 2;
    }
    return m;
}

int main() {
    int n;
    scanf ("%d", &n);
    printf("bitrev(%d)=%d\n", n, bitrev(n));
    return 0;
}
```

#### **Problem 4: job scheduling** (20 points)

A couple of scientists want to run jobs on a super computer. Jobs on this super computer are executed one after another (serially). Each scientist knows in advance how long his job will run. They are all eager for their results, and therefore they do not want to wait too long for their results. The system administrator of the computer decides to run the jobs in an order that minimizes the *average waiting time*. The waiting time for a scientist is the time between supplying the job (for all jobs, this is time t=0) and the time that his job finished.

As a simple example, consider the case with only two jobs. The expected execution time for job 1 is 10 hours, while the execution time of job 2 is 4 hours. If we run job 1 first, then the total waiting time is 10 + (4+10) = 24 hours (and hence the average waiting time is 12 hours). However, if we run job 2 first, then the total waiting time will be 4 + (10+4)=18 hours (and the average waiting time would be 9 hours). Therefore, the system administrator decides to run job 2 first.

Write a program that computes the order in which the jobs should be run in order to minimize the average waiting time. The first line of the input is an integer n, the number of jobs. You may assume that 1 < n < 100. The rest of the input consists of n lines, which consist of the running times of the n jobs (in hours). You may assume that all running times are integer numbers less than 100. Of course, the first running time belongs to job 1, the second to job 2, and so on. The output should consist of 1 line, which is the order in which the jobs need to be run. The job numbers are separated by a comma. The output line must end with a newline (i.e the '\n' character). Note, that the first job has job number 1. If two (or more) jobs have the same running time, then these jobs must be ordered based on their job number (smallest number first).

```
for (i = 0; i < n; i++) {
    if (job[i] == min) {
        if (cnt > 0) {
            printf(",");
        }
        printf("%d", i+1);
        job[i] = 9999;
        cnt++;
        }
    }
    printf("\n");
    return 0;
}
```

#### **Problem 5: ordered bit strings using recursion** (20 points)

In this problem we consider bit strings, i.e. sequences consisting solely of zeros and ones. Consider the following example bit string: 010011. This string satisfies the property that *each* prefix (i.e. leading substring) of the string (including the string itself) contains at most as many ones as it contains zeros. It is easy to verify this property for all prefixes of the example string: 0, 01, 010, 0100, 01001, and 010011. An example of a string for which this property does not hold is the string 1100, since each of the prefixes 1, 11, 110 contain more ones than zeros.

The input of this problem is an integer n, where  $0 < n \le 10$ . The output must be all bit strings of length 2n, containing n zeros and n ones, that satisfy the above mentioned property. Moreover, the strings must be ordered in lexicographical (dictionary sorted) order. Each string must be printed on a separate line without spaces or tabs (see output examples).

The following incomplete code fragment is available in Justitia. Download it and complete the code. You are asked to implement the body of the function <code>showBitStrings</code>. This function should call a recursive helper function (with suitably chosen parameters/arguments) that solves the problem. You are not allowed to make changes in the main function.

```
void bitstrings(int zeros, int ones, int len, int bits[]) {
 if (zeros == 0 && ones == 0) {
    int i;
    for (i=0; i < len; i++) {
      printf("%d", bits[i]);
   printf("\n");
  } else {
    if (zeros > 0) {
      bits[len] = 0;
      bitstrings(zeros-1, ones, len+1, bits);
    if (zeros < ones) {
      bits[len] = 1;
      bitstrings(zeros, ones-1, len+1, bits);
  }
}
void showBitStrings(int n) {
 int bits[20];
 bitstrings(n, n, 0, bits);
int main() {
 int n;
 scanf ("%d", &n);
 showBitStrings(n);
 return 0;
```